

## OPTIMAL SCHEDULING OF THE “INVIGILATION” PROBLEM

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We define the *invigilation* problem as follows: A university department offers  $m$  modules per term, and there are  $k$  available invigilators for their exams. Depending on the enrollment, each module  $j$  requires  $w_j$  invigilators,  $w \in \mathbb{N}^m$ . Each invigilator  $i$  is only available for the exams of certain modules and cannot be assigned more than  $v_i$  invigilations per term,  $v \in \mathbb{N}^k$ . The availability of invigilators to modules can be described as a boolean  $k \times m$  matrix  $E$ , with  $E_{ij} = 1$ , if person  $i$  can *potentially* be present at the exam of module  $j$ , and  $E_{ij} = 0$  otherwise.

This model problem represents a number of scheduling applications from diverse fields such as workplace staffing, radio/tv broadcasts, marketing, production processes, etc.

An optimal scheduling solution for the “invigilation” problem is a boolean matrix  $X \in \mathbb{N}^{k \times m}$ , with  $X_{ij} = 1$ , if invigilator  $i$  is assigned to module  $j$  and  $X_{ij} = 0$  otherwise. The solution should

$$\text{minimize } \sum_{j=1}^m X_{ij}, \quad \text{for } i = 1, 2, \dots, k,$$

subject to the following constraints:

$$\sum_{j=1}^m X_{ij} \leq v_i, \quad \text{for } i = 1, 2, \dots, k \quad \text{and} \quad \sum_{i=1}^k X_{ij} \geq w_j, \quad \text{for } j = 1, 2, \dots, m.$$

We propose various numerical and heuristic algorithms for the solution of the “invigilation” problem and we investigate their performance on a number of test problems.